







## An algebra of variant loci

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#### Introduction

Variant calling on NGS data often entails filtering samples against each other to e.g.

- detect de-novo mutations (child vs. parents, tumor vs. normal),
- eliminate sequencing artifacts.

This can be formulated as set operations, e.g.

$$V_A \setminus (V_B \cup V_C)$$

with  $V_A, V_B$  and  $V_C$  being the true variant loci of sample A, B and C. Of course, these sets are unknown. The state of the art is to call variants of each sample, and perform set-based filtering afterwards.

This gives rise to three problems:

**Insufficient evidence problem** The filtering fails if the coverage is too low.

N+1 problem Calling samples in groups helps with the insufficient evidence problem. But later addition of new samples leads to redundant calculations.

**FDR problem** The obtained variant qualities do not reflect the filtering. This makes controlling the false discovery rate (FDR) difficult.

Apart from specialized solutions for Tumor/Normal pairs, no solution for generic filtering scenarios exists that solves all problems.

We present **Algebraic Variant Calling**, an approach to incorporate the filtering into the calling model. Algebraic Variant Calling solves the insufficient evidence problem and provides intuitive FDR control. In the ALgebraic PArallel CAller (**ALPACA**), we combine Algebraic Variant Calling with a BCF based approach to solve the N+1 problem.

#### Algebra of variant loci

For a finite set of samples  $S=\{s_1,s_2,\dots\}$  with variant loci  $V_S=\{V_{s_1},V_{s_2},\dots\}$ , we define the algebra

$$\mathcal{A}_S = \left(2_S^V \setminus \emptyset, \cup, \setminus, \left(\bigotimes^k\right)_{k \in \mathbb{N}}\right)$$

with the classic set operations union  $\cup$  and difference  $\setminus$  and a k-relaxed intersection  $\bigotimes^k$ . The k-relaxed intersection  $\bigotimes^k_{s \in S'} V_s$  for subset  $S' \subseteq S$  with  $|S'| \ge k$  is the set of variant loci common to at least k of the samples in S'.

This allows all kinds of filtering scenarios, e.g.

ullet Call all variants in any of the samples a,b,c:

$$V_a \cup V_b \cup V_c$$

ullet Call somatic mutations in e.g. a tumor sample t compared to a healthy normal sample n:

$$V_t \setminus V_n$$

ullet Call de-novo mutations in a child sample c compared to its parents f,m:

$$V_c \setminus (V_f \cup V_m)$$

ullet Call somatic mutations in a group of tumors t,t' compared to their normals n,n':

$$(V_t \cup V_{t'}) \setminus (V_n \cup V_{n'})$$

• Do the same in a paired way:

$$(V_t \setminus V_n) \cup (V_{t'} \setminus V_{n'})$$

ullet Call all variants that are recurrent in at least 3 of the samples a,b,c,d,e:

$$\bigotimes_{s \in \{a,b,c,d,e\}}^{3} V$$

### Algebraic variant calling

For any scenario  $\phi \in \mathcal{A}_S$ , we calculate the posterior probability for the null hypothesis  $i \notin \phi$  for each locus i. If  $\phi = \bigcup_{s \in S' \subset S} V_s$ , we calculate

$$\Pr(i \not\in \phi \mid D_{S,i}) := \Pr(M = 0 \mid D_{S',i})$$

e.g. in the usual Bayesian way (dePristo et al. 2011, Li 2010). If  $\phi=\phi_1\setminus\phi_2$ , we write

$$\Pr(i \not\in \phi \mid D_{S,i}) := 1 - \Pr(i \in \phi_1 \mid D_{S,i}) \cdot \Pr(i \not\in \phi_2 \mid D_{S,i}),$$

and  $\phi = \phi_1 \cup \phi_2$  leads to

$$\Pr(i \not\in \phi \mid D_{S,i}) := \Pr(i \not\in \phi_1 \mid D_{S,i}) \cdot \Pr(i \not\in \phi_2 \mid D_{S,i}).$$

For the k-relaxed intersection  $\phi = \bigotimes_{i=1,2,\dots}^k \phi_i$  we can calculate  $\Pr(i \notin \phi \mid D_{S,i})$  with dynamic programming.

Finally, we can approximate  $\phi$  as

$$\phi_{\alpha}^* := \{i \mid \forall i = 1, 2, \dots, n : \Pr(i \in \phi | D_{S,i}) \le \alpha \}.$$

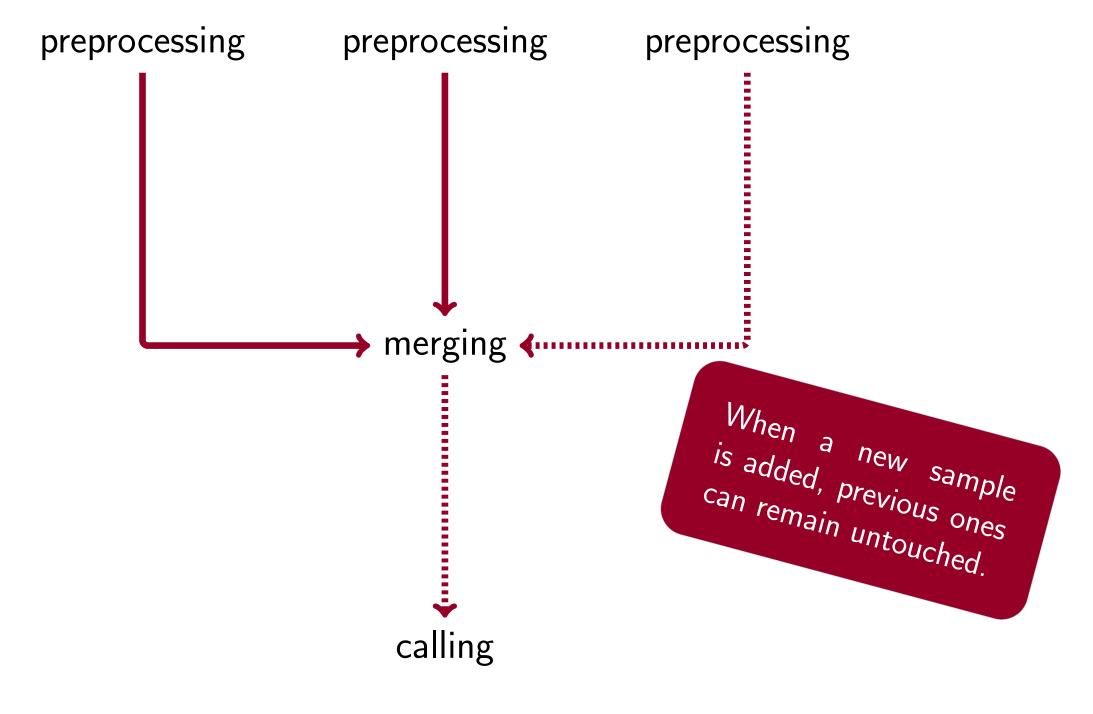
# Solving the N+1 problem

The probability  $\Pr(M=0|D_{S,i})$  is calculated from per-sample genotype likelihoods

$$\Pr(D_{s,i}|G=g)$$

with G being the random variable denoting possible genotypes.  $D_{s,i}$  is the pileup of read bases of sample s at locus i. The likelihoods are independent of the query formula  $\phi$ , hence:

- Genotype likelihoods for all covered loci can be preprocessed into persample BCF files.
- Sample BCF files can be merged into a global BCF, keeping only loci with any non-reference maximum likelihood genotype.
- ullet Calling with different scenarios  $\phi \in \mathcal{A}_S$  becomes a matter of seconds.



#### **Controlling FDR**

FDR can be controlled to not exceed  $lpha^*$  by setting the threshold

$$\alpha = \max\{\alpha' \in [0, \alpha^*] \mid \overline{FDR}_{\alpha'} \le \alpha^*\}$$

with

$$\overline{FDR}_{\alpha} = \frac{1}{|\phi_{\alpha}^*|} \sum_{i \in \phi^*} \Pr(i \not\in \phi | D_{S,i}).$$

Since posterior probabilities reflect the given query, controlling FDR becomes easy.